NEW ESTIMATION TECHNIQUE FOR A CLASS OF CHAOTIC SIGNALS

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ABSTRACT
We propose a new technique for the estimation of chaotic signals in the presence of additive noise. The new method uses statistical measures to restrict the space over which the signals are received. The proposed technique is applicable to a large variety of chaotic signals and has good performance indicated by the low estimation error bias and variance. The complexity of the algorithm is shown to be low.

1. INTRODUCTION
Chaotic systems have very desirable properties for signal analysis and synthesis. Although they produce random noise like signals, they are actually the output of deterministic systems. They also have a very wide frequency spectrum that make it difficult to intercept and detect. As a result they are very promising candidates for application in communication systems. The key feature of these systems is their sensitive dependence on initial conditions.

The important issue in using a chaotic signal for communication is the existence of a ‘good’ receiver and the ease of its implementation. Maximum likelihood estimation algorithms are known only for a very limited class of chaotic signals. Characteristics of asymptotic maximum likelihood estimators for chaotic signals are derived in [1] and a specific maximum likelihood estimator for the tent map is given in [4].

In this paper we propose a new technique for the estimation of a wide range of chaotic signals. The proposed algorithm is not restricted to AWGN channels but could also be applied to any kind of additive noise channels. We show that the new algorithm gives lower error variances than existing algorithms. The proposed algorithm is not restricted to piecewise linear or unimodal maps.

2. SIGNAL AND CHANNEL MODELS
We consider the estimation of signals that are corrupted by additive white noise. The channel model under consideration is given by

\[ y_n = x_n + z_n, \]

where \( x_n \) is the transmitted signal, \( y_n \) is the received signal, \( z_n \) is the noise component and \( n \) represents the time index. The transmitted signal sequence \( x_n \) is generated as the output of a chaotic map. In this paper we restrict ourselves to one-dimensional maps as specified below.

\[ x_n = f(x_{n-1}) \]

where \( f(\cdot) \) is a non linear function. However, we do not restrict ourselves to unimodal or piecewise linear maps, which is a requirement on other algorithms like those proposed in [2]. Some examples of chaotic maps are given in Fig 1. The additive noise is not assumed to have Gaussian distribution in the estimation algorithm.

3. ESTIMATION TECHNIQUE
Our basic problem is to find an estimate of the transmitted signal \( x_n \) for \( n = 1, 2, \ldots, N \). This problem reduces to accurately estimating the initial
transmitted signal (or initial condition of the map). However, using steepest gradient methods to directly estimate the initial condition is not possible because of the fractal nature of the estimation error [3].

Instead [4] suggests a scheme wherein the itinerary of the signal is first determined and this information is used to find the initial condition. The itinerary of a chaotic sequence $x_1, x_2, ..., x_N$ generated from a unimodal map is defined by

$$p_n = \begin{cases} 0 & \text{if } 0 \leq x_n < c \\ 1 & \text{if } c \leq x_n < 1 \end{cases}$$

for $n = 1, 2, ..., N-1$ and $c$ is the value of $x$ for which $f(x)$ is maximum. This definition can be easily modified to include other maps like those in Fig 1. We will illustrate our algorithm for the tent map which is given by

$$f(x) = \beta - 1 - \beta(x),$$

where $1 < \beta \leq 2$.

Assume initially that we are operating in the high SNR region. The main idea is that if the itinerary of the transmitted signals is known then we can effectively use a steepest descent algorithm to estimate the final condition. Using this estimate of the final condition and the itinerary we can find an estimate of the initial condition. Our contention is that with a large probability the itineraries of the received signal and transmitted signal are almost equal. So we propose that we use a test to search over a subspace of possible itineraries.

For the case of tent map, the itinerary of the signal sequence is the same as the sign of the transmitted signal sequence. We include all possible values for the $i^{th}$ component of the itinerary if $|y_{i-1} - \alpha_1| \leq \epsilon$ or $|y_{i-1} - \alpha_2| \leq \epsilon$ where $\alpha_1$ and $\alpha_2$ are the zeros of $f()$. Otherwise we take the itinerary to be equal to the sign of the received signal. Using this we form a list of most likely itineraries.

It can be easily shown that the probability of not including the correct itinerary of the transmitted signal sequence in the list of possible itineraries can be approximated by

$$2 \Pr (\alpha - \epsilon < x_n < \alpha + \epsilon) \Pr (z_n > \epsilon)$$

which reduces to $2\epsilon Q(\frac{\epsilon}{\sigma})$ under the assumption that $z$ is $N(0, \sigma^2)$ and $x_n$ is uniformly distributed between -1 and 1, which is an invariant density for the tent map.

The plot of this probability for different $\epsilon$ is given in Fig 2 for Gaussian noise case at an SNR of 27dB. It is clear from the figure that for any value of $\epsilon$ the probability that our subset of itineraries will not include the correct itinerary is below $10^{-3}$. The size of the search set depends on the value of $\epsilon$ and can be used to get a tradeoff between complexity and performance.

![Figure 2: Error in not including the correct itinerary in our subset of possible itineraries for tent map, SNR of 20dB](image)
only searching over a subset of possible iternaries, we call our new estimator as subspace based estimator.

The new subspace algorithm is summarized below

- Choose an $\epsilon$ based on knowledge of SNR and amount of computation allowed
- Form the subset of possible iternaries based on whether or not $|y_{k-1} - \alpha_1| \leq \epsilon$ or $|y_{k-1} - \alpha_2| \leq \epsilon$ where $\alpha_1$ and $\alpha_2$ are the zeros of $f()$.
- For each itenary in selected subset estimate the final value $x_N$ using a grid search.
- Choose the itenary and final value that gives the lowest estimation error

4. RESULTS

The plot of the estimation error versus SNR is shown in Fig 4 for the tent map and in Fig 3 for the logistic map for an AWGN channel. From this figure it is clear that our error is lower than the Cramer Rao bound and consequently biased. For the logistic map we have indicated the performance of the subspace estimator for two different values of resolution of grid for estimating the final value. The subspace of possible iternaries over which we search is maintained the same in both the cases. As can be expected the higher resolution search achieves Cramer-Rao bound whereas the lower resolution search performs badly especially at high SNR’s. The plot of the estimator bias is shown in Fig 5. It is clear from the figure that the estimator bias is quite small especially for high SNR’s. It is our conjecture that at high SNR’s the estimator is essentially unbiased and achieves the Cramer-Rao bound.

5. CONCLUSIONS

We have introduced a new technique for the estimation of a wide range of chaotic signals. The new technique involves enumeration of a set of most likely iternaries and then jointly estimating the final condition and itenary of the map. This subspace based search algorithm has low implementation complexity and low estimator error variance. It is applicable to all kinds of additive noise channels and not necessarily Gaussian channels. It can also be easily extended to the case in which some known function of the chaotic signal is observed in noise. This will encompass the case of slowly fading or block fading channels. In the case of multiple chaotic signals using the same physical medium and interfering with one another different methods will be needed to achieve single user performance.
6. REFERENCES


